

Lists and Recursion and Trees

Oh My!

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T.P.S. REPORT

COVER SHEET

Prepared By: _____ Date: _____

Device/Program Type: _____

Product Code: _____ Customer: _____

Vendor: _____

Due Date: _____ Data Loss: _____

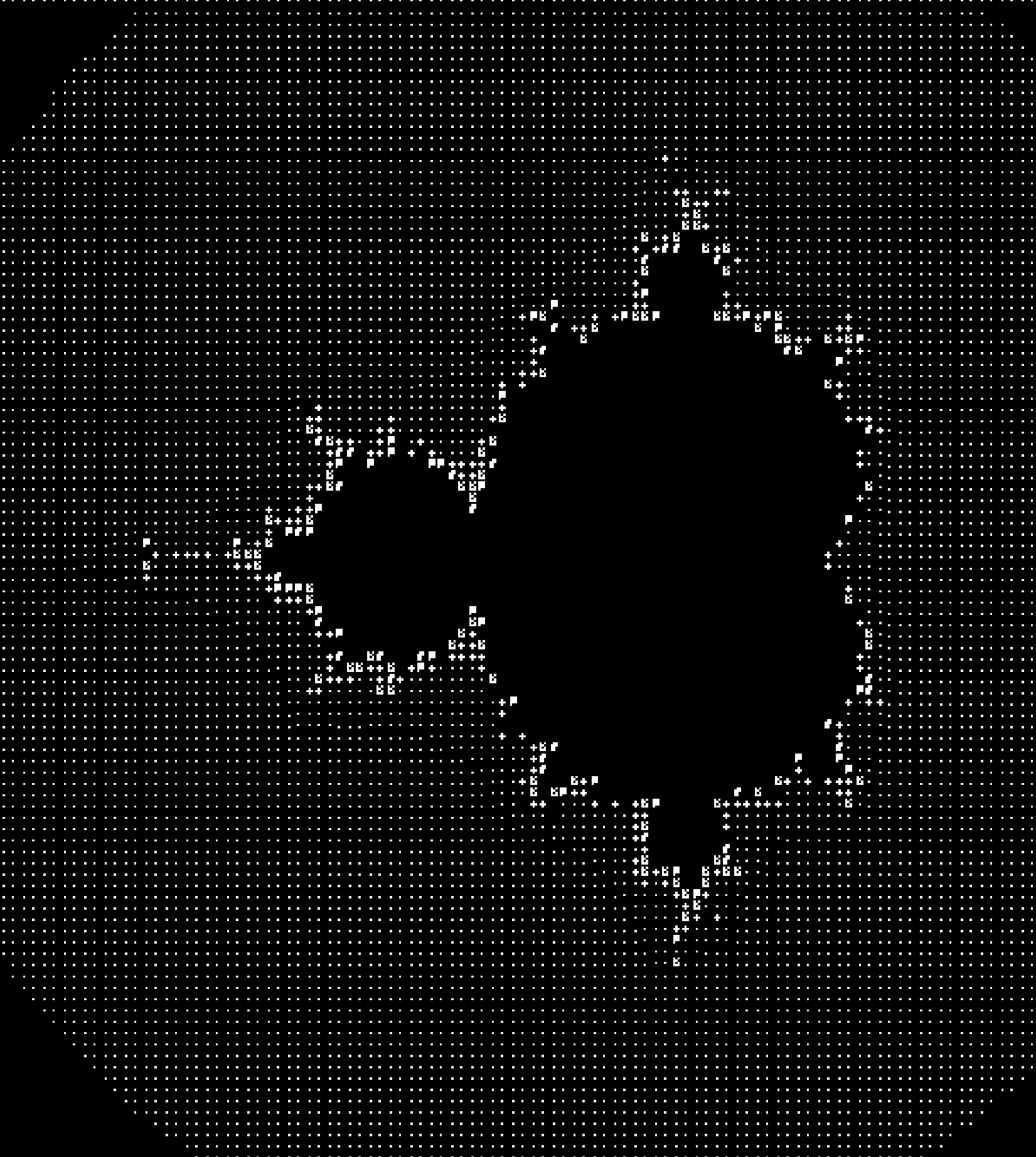
Test Date: _____ Target Run Date: _____

Program Run Time: _____ Reference Guide: _____

Program Language: _____ Number of Error Messages: _____

Comments: _____

C O N F I D E N T I A L



Better, Faster TPS Reports

New!

Reach Outside the Current Row

Better, Faster TPS Reports

- **Windowing Function**
 - Operates on a window
 - Returns a value for each row
 - Calculates value from the rows in the window

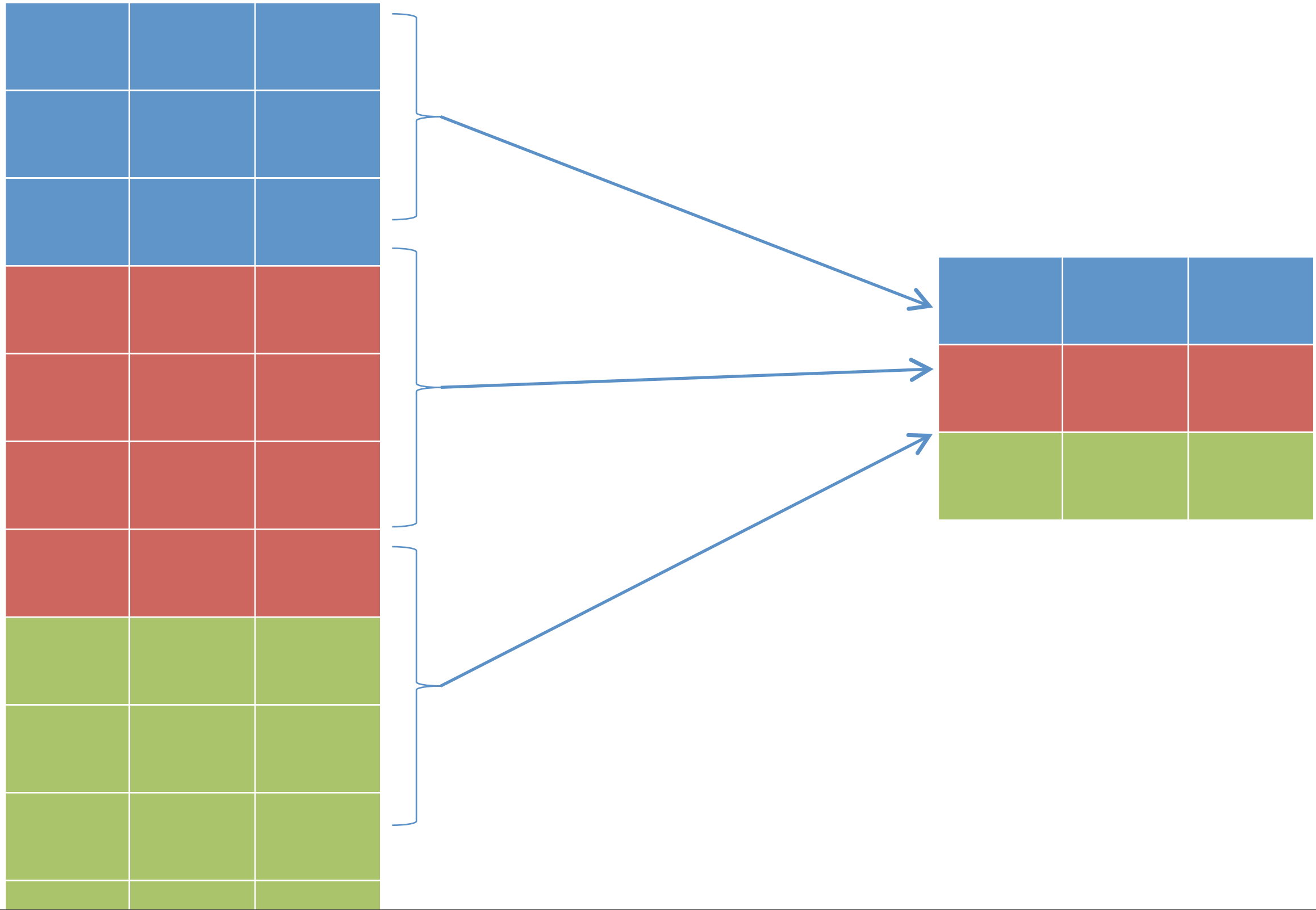
Better, Faster TPS Reports

- You can use...
 - New window functions
 - Existing aggregate functions
 - User-defined window functions
 - User-defined aggregate functions

Better, Faster TPS Reports

[Aggregates]

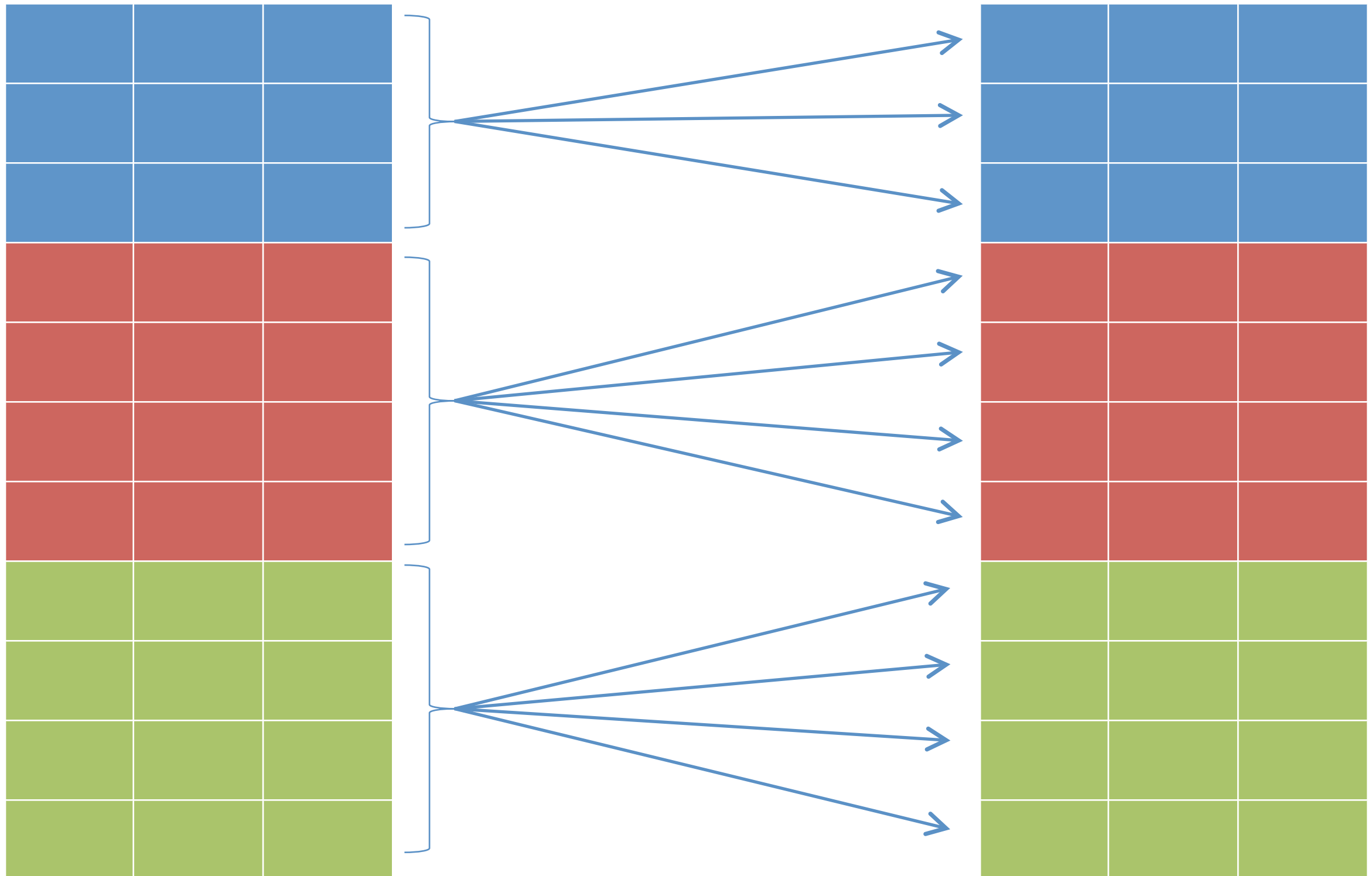
```
SELECT key, SUM(val) FROM tbl GROUP BY key;
```



Better, Faster TPS Reports

[Windowing Functions]

SELECT key, SUM(val) OVER (PARTITION BY key) FROM tbl;



ROW_NUMBER (Before)

```
SELECT
    e1.empno,
    e1.depname,
    e1.salary,
    count(*) AS row_number
FROM
    empsalary e1
JOIN
    empsalary e2
    ON (e1.empno < e2.empno)
GROUP BY e1.empno, e1.depname, e1.salary
ORDER BY e1.empno DESC;
```

ROW_NUMBER (Before)

OOPS!

empno	depname	salary	row_number
8	develop	6000	1
6	sales	5500	2
11	develop	5200	4
10	develop	5200	4
1	sales	5000	5
3	sales	4800	7
4	sales	4800	7
9	develop	4500	8
7	develop	4200	9
2	personnel	3900	10
5	personnel	3500	11

(11 rows)

ROW_NUMBER (After)

```
SELECT
    empno,
    depname,
    salary,
    row_number() OVER (
        ORDER BY salary DESC NULLS LAST
    )
FROM
    empsalary
ORDER BY salary DESC;
```

ROW_NUMBER (After)

Yippee!

empno	depname	salary	row_number
8	develop	6000	1
6	sales	5500	2
10	develop	5200	3
11	develop	5200	4
1	sales	5000	5
3	sales	4800	6
4	sales	4800	7
9	develop	4500	8
7	develop	4200	9
2	personnel	3900	10
5	personnel	3500	11

(11 rows)

More Ranking

```
SELECT
    empno,
    depname,
    salary,
    row_number() OVER (
        ORDER BY salary DESC NULLS LAST
    ),
    rank() OVER (
        ORDER BY salary DESC NULLS LAST
    ),
    dense_rank() OVER (
        ORDER BY salary DESC NULLS LAST
    )
FROM
    empsalary
ORDER BY salary DESC;
```

More Ranking

empno	depname	salary	row_number	rank	dense_rank
8	develop	6000	1	1	1
6	sales	5500	2	2	2
10	develop	5200	3	3	3
11	develop	5200	4	3	3
1	sales	5000	5	5	4
3	sales	4800	6	6	5
4	sales	4800	7	6	5
9	develop	4500	8	8	6
7	develop	4200	9	9	7
2	personnel	3900	10	10	8
5	personnel	3500	11	11	9

(11 rows)

Built-in Windowing Functions

- `row_number()`
- `rank()`
- `dense_rank()`
- `percent_rank()`
- `cume_dist()`
- `ntile()`
- `lag()`
- `lead()`
- `first_value()`
- `last_value()`
- `nth_value()`

row_number()

- Returns number of the current row

```
SELECT val, row_number() OVER (ORDER BY val DESC) FROM tbl;
```

val	row_number()
5	1
5	2
3	3
1	4

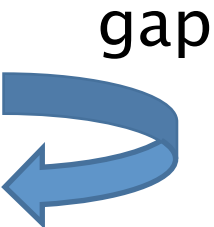
Note: row_number() always incremented values independent of frame

rank()

- Returns rank of the current row **with** gap

```
SELECT val, rank() OVER (ORDER BY val DESC) FROM tbl;
```

val	rank()
5	1
5	1
3	3
1	4




Note: rank() OVER(*empty*) returns 1 for all rows, since all rows are peers to each other

dense_rank()

- Returns rank of the current row **without** gap

```
SELECT val, dense_rank() OVER (ORDER BY val DESC) FROM tbl;
```

val	dense_rank()
5	1
5	1
3	2
1	3



no gap

Note: `dense_rank() OVER(*empty*)` returns 1 for all rows, since all rows are peers to each other

percent_rank()

- Returns relative rank; $(\text{rank}() - 1) / (\text{total row} - 1)$

`SELECT val, percent_rank() OVER (ORDER BY val DESC) FROM tbl;`

val	percent_rank()
5	0
5	0
3	0.6666666666666667
1	1

values are always between 0 and 1 inclusive.

Note: `percent_rank() OVER(*empty*)` returns 0 for all rows, since all rows are peers to each other

cume_dist()

- Returns relative rank; (# of preced. or peers) / (total row)

```
SELECT val, cume_dist() OVER (ORDER BY val DESC) FROM tbl;
```

val	cume_dist()	
5	0.5	= 2 / 4
5	0.5	= 2 / 4
3	0.75	= 3 / 4
1	1	= 4 / 4

The result can be emulated by

“count(*) OVER (ORDER BY val DESC) / count(*) OVER ()”

Note: cume_dist() OVER(*empty*) returns 1 for all rows, since all rows are peers to each other

ntile()

- Returns dividing bucket number

```
SELECT val, ntile(3) OVER (ORDER BY val DESC) FROM tbl;
```

val	ntile(3)
5	1
5	1
3	2
1	3

→ $4 \% 3 = 1$

The results are the divided positions, but if there's remainder add row from the head

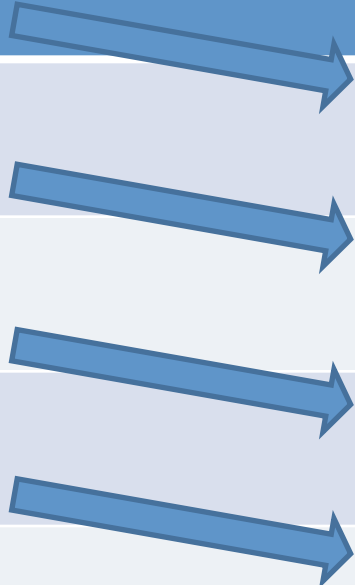
Note: `ntile() OVER (*empty*)` returns same values as above, since `ntile()` doesn't care the frame but works against the partition

lag()

- Returns value of row above

```
SELECT val, lag(val) OVER (ORDER BY val DESC) FROM tbl;
```

val	lag(val)
5	NULL
5	5
3	5
1	3



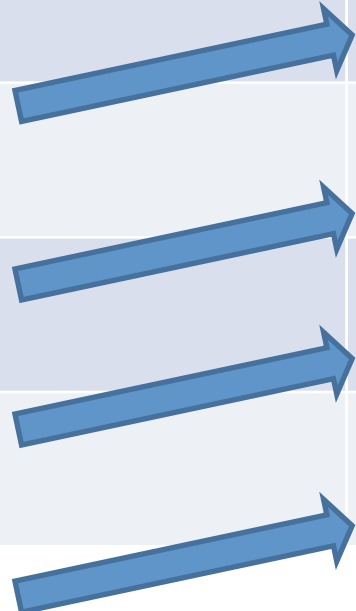
Note: lag() only acts on a partition.

lead()

- Returns value of the row below

```
SELECT val, lead(val) OVER (ORDER BY val DESC) FROM tbl;
```

val	lead(val)
5	5
5	3
3	1
1	NULL



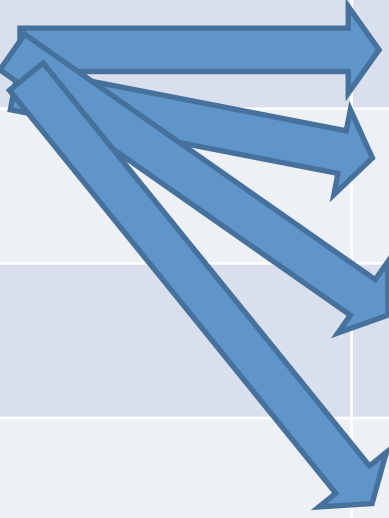
Note: lead() acts against a partition.

first_value()

- Returns the first value of the frame

```
SELECT val, first_value(val) OVER (ORDER BY val DESC) FROM tbl;
```

val	first_value(val)
5	5
5	5
3	5
1	5

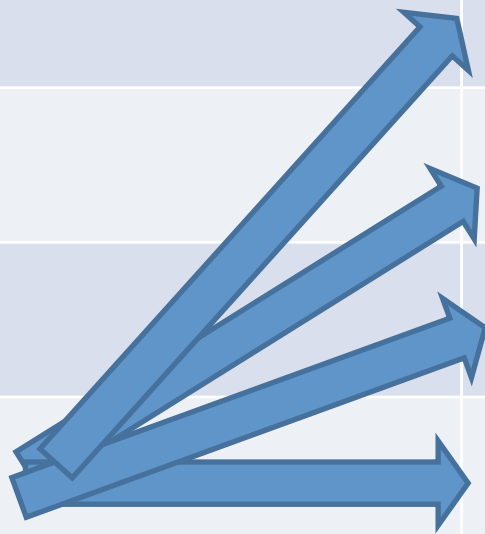


last_value()

- Returns the last value of the frame

```
SELECT val, last_value(val) OVER  
(ORDER BY val DESC ROWS BETWEEN UNBOUNDED PRECEDING  
AND UNBOUNDED FOLLOWING) FROM tbl;
```

val	last_value(val)
5	1
5	1
3	1
1	1



Note: frame clause is necessary since you have a frame between the first row and the current row by only the order clause

nth_value()

- Returns the n-th value of the frame

```
SELECT val, nth_value(val, val) OVER  
(ORDER BY val DESC ROWS BETWEEN UNBOUNDED PRECEDING  
AND UNBOUNDED FOLLOWING) FROM tbl;
```

val	nth_value(val, val)
5	NULL
5	NULL
3	3
1	5

Note: frame clause is necessary since you have a frame between the first row and the current row by only the order clause

aggregates(all peers)

- Returns the same values along the frame

```
SELECT val, sum(val) OVER () FROM tbl;
```

val	sum(val)
5	14
5	14
3	14
1	14

Note: all rows are the peers to each other

cumulative aggregates

- Returns different values along the frame

```
SELECT val, sum(val) OVER (ORDER BY val DESC) FROM tbl;
```

val	sum(val)
5	10
5	10
3	13
1	14

Note: row#1 and row#2 return the same value since they are the peers.
the result of row#3 is sum(val of row#1...#3)



recursion

Search

[Advanced Search](#)Search: the web pages from AustraliaWeb [+ Show options...](#)

Results 1 - 10 of about 715,000 for recursion [definition]. (0.08 seconds)

Did you mean: [recursion](#)[Recursion - Wikipedia, the free encyclopedia](#)

A visual form of **recursion** known as the Droste effect. The woman in this image is holding an object which contains a smaller image of her holding the same ...

en.wikipedia.org/wiki/Recursion - [Cached](#) - [Similar](#)

[Recursion \(computer science\) - Wikipedia, the free encyclopedia](#)

Recursion in computer science is a method where the solution to a problem depends on solutions to smaller instances of the same problem. ...

[en.wikipedia.org/wiki/Recursion_\(computer_science\)](http://en.wikipedia.org/wiki/Recursion_(computer_science)) - [Cached](#) - [Similar](#)

[+ Show more results from en.wikipedia.org](#)[Recursion -- from Wolfram MathWorld](#)

A **recursive** process is one in which objects are defined in terms of other objects of the same type. Using some sort of recurrence relation, the entire class ...

mathworld.wolfram.com › ... › [Algorithms](#) › [Recursion](#) - [Cached](#) - [Similar](#)

[recursion](#)

Definition of **recursion**, possibly with links to more information and implementations.

www.itl.nist.gov/div897/sqg/dads/HTML/recursion.html - [Cached](#) - [Similar](#)

[Did you mean recursion?](#)

23 Jul 2009 ... Bloodwine, on 07/23/2009, -8/+206Yo Dawg I herd you like **recursion** so we put **recursion** in yo **recursion** so you can repeat while you repeat ...

Find: Highlight all Match case

Done

Generate Points

```
WITH RECURSIVE x(i)
AS (
  VALUES (0)
UNION ALL
  SELECT i + 1
  FROM x
  WHERE i < 101
),
```

Generate Points

```
Z(Ix, Iy, Cx, Cy, X, Y, I)
AS (
    SELECT Ix, Iy,
           X::float, Y::float,
           X::float, Y::float,
           0
    FROM
```

Generate Points

```
(SELECT -2.2 + 0.031 * i, i
FROM x) AS xgen(x,ix)
CROSS JOIN
(SELECT -1.5 + 0.031 * i, i
FROM x) AS ygen(y,iy)
```


Generate Points

UNION ALL

Generate Points

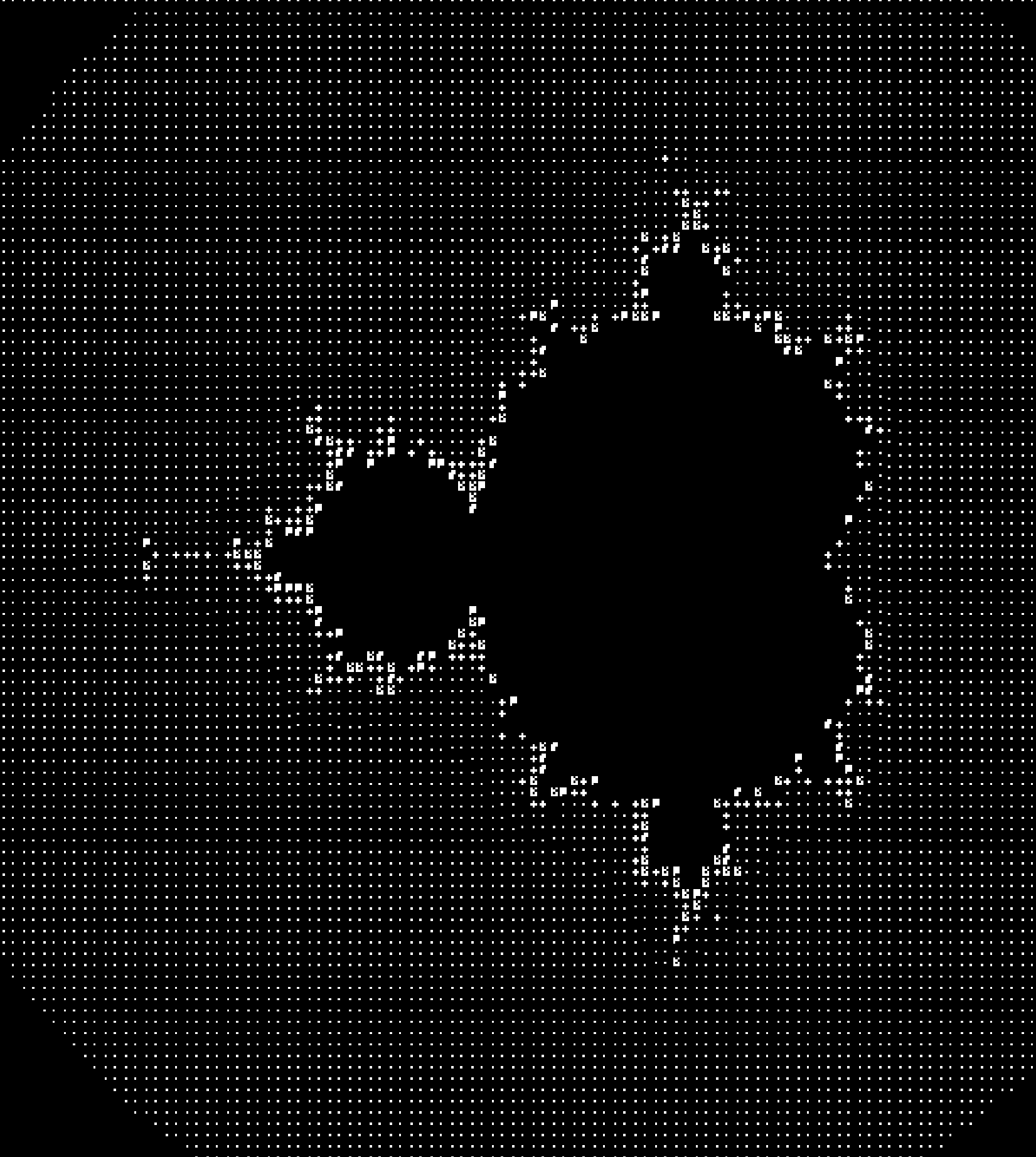
```
SELECT
    Ix, Iy, Cx, Cy,
    X * X - Y * Y + Cx AS X,
    Y * X * 2 + Cy,
    I + 1
FROM Z
WHERE X * X + Y * Y < 16.0
AND I < 27
),
```

Choose Some

```
Zt (Ix, Iy, I) AS (  
    SELECT Ix, Iy, MAX(I) AS I  
    FROM Z  
    GROUP BY Iy, Ix  
    ORDER BY Iy, Ix  
)
```

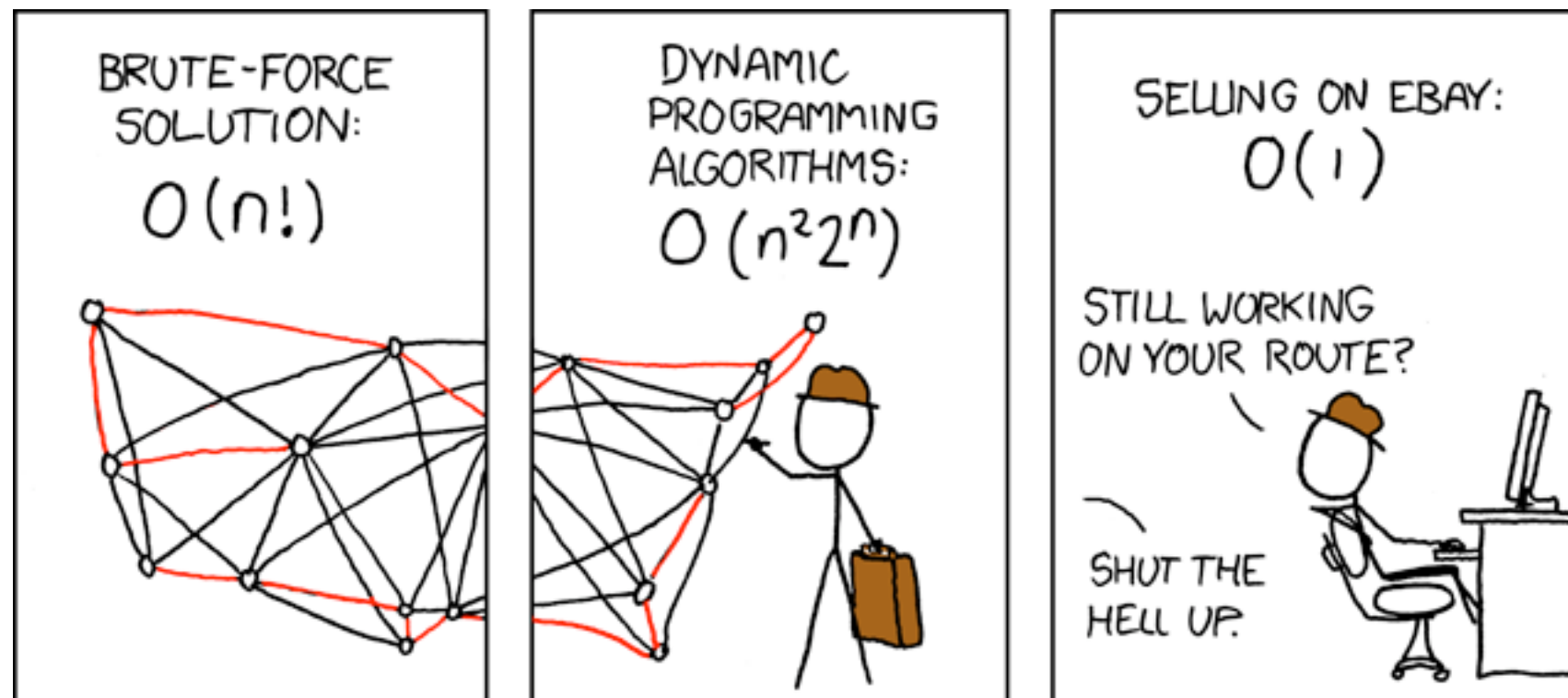
Display Them

```
SELECT array_to_string(  
    array_agg(  
        SUBSTRING(  
            ' . , , , - - - - - + + + + + % % % % @ @ @ @ ##### ' ,  
            GREATEST(I, 1)  
        ), ''  
    )  
FROM zt  
GROUP BY Iy  
ORDER BY Iy;
```



Travelling Salesman Problem

Given a number of cities and the costs of travelling from any city to any other city, what is the least-cost round-trip route that visits each city exactly once and then returns to the starting city?



TSP Schema

```
CREATE TABLE pairs (  
    from_city TEXT NOT NULL,  
    to_city TEXT NOT NULL,  
    distance INTEGER NOT NULL,  
    PRIMARY KEY(from_city, to_city),  
    CHECK (from_city < to_city)  
);
```

TSP Data

```
INSERT INTO pairs
```

```
VALUES
```

```
    ('Bari', 'Bologna', 672),  
    ('Bari', 'Bolzano', 939),  
    ('Bari', 'Firenze', 723),  
    ('Bari', 'Genova', 944),  
    ('Bari', 'Milan', 881),  
    ('Bari', 'Napoli', 257),  
    ('Bari', 'Palermo', 708),  
    ('Bari', 'Reggio Calabria', 464),  
    ....
```


TSP Program:

Symmetric Setup

```
WITH RECURSIVE both_ways (  
    from_city,  
    to_city,  
    distance  
)          /* Working Table */  
AS (  
    SELECT  
        from_city,  
        to_city,  
        distance  
    FROM  
        pairs  
    UNION ALL  
    SELECT  
        to_city AS "from_city",  
        from_city AS "to_city",  
        distance  
    FROM  
        pairs  
) ,
```

TSP Program:

Symmetric Setup

```
WITH RECURSIVE both_ways(  
    from_city,  
    to_city,  
    distance  
)  
AS (/* Distances One Way */  
    SELECT  
        from_city,  
        to_city,  
        distance  
    FROM  
        pairs  
UNION ALL  
    SELECT  
        to_city AS "from_city",  
        from_city AS "to_city",  
        distance  
    FROM  
        pairs  
) ,
```

TSP Program:

Symmetric Setup

```
WITH RECURSIVE both_ways(  
    from_city,  
    to_city,  
    distance  
)  
AS (  
    SELECT  
        from_city,  
        to_city,  
        distance  
    FROM  
        pairs  
UNION ALL /* Distances Other Way */  
    SELECT  
        to_city AS "from_city",  
        from_city AS "to_city",  
        distance  
    FROM  
        pairs  
) ,
```

TSP Program:

Path Initialization Step

```
paths (  
    from_city,  
    to_city,  
    distance,  
    path  
)  
AS (  
    SELECT  
        from_city,  
        to_city,  
        distance,  
        ARRAY[from_city] AS "path"  
    FROM  
        both_ways b1  
    WHERE  
        b1.from_city = 'Roma'  
UNION ALL
```

TSP Program:

Path Recursion Step

```
SELECT
    b2.from_city,
    b2.to_city,
    p.distance + b2.distance,
    p.path || b2.from_city
FROM
    both_ways b2
JOIN
    paths p
ON (
    p.to_city = b2.from_city
AND
    b2.from_city <> ALL (p.path[
        2:array_upper(p.path,1)
    ]) /* Prevent re-tracing */
AND
    array_upper(p.path,1) < 6
)
)
```

TSP Program:

Timely Termination Step

```
SELECT
    b2.from_city,
    b2.to_city,
    p.distance + b2.distance,
    p.path || b2.from_city
FROM
    both_ways b2
JOIN
    paths p
ON (
    p.to_city = b2.from_city
AND
    b2.from_city <> ALL (p.path[
        2:array_upper(p.path,1)
    ]) /* Prevent re-tracing */
AND
    array_upper(p.path,1) < 6 /* Timely Termination */
)
)
```

TSP Program:

Filter and Display

```
SELECT
    path || to_city AS "path",
    distance
FROM
    paths
WHERE
    to_city = 'Roma'
AND
    ARRAY['Milan', 'Firenze', 'Napoli'] <@ path
ORDER BY distance, path
LIMIT 1;
```

TSP Program:

Filter and Display

```
 davidfetter@tsp=# \i travelling_salesman.sql
                path                | distance
-----+-----
 {Roma, Firenze, Milan, Napoli, Roma} |      1553
(1 row)
```

Time: 11679.503 ms

Who Posts Most?

Who

```
CREATE TABLE forum_users (  
    user_name TEXT NOT NULL,  
    CHECK(user_name = trim(user_name)),  
    user_id SERIAL UNIQUE  
);  
  
CREATE UNIQUE INDEX forum_user_user_name_unique  
    ON forum_users(lower(user_name));  
  
INSERT INTO forum_users (user_name)  
VALUES  
    ('Tom Lane'), ('Robert Haas'), ('Alvaro Herrera'), ('Dave Page'),  
    ('Heikki Linnakangas'), ('Magnus Hagander'), ('Gregory Stark'),  
    ('Josh Berkus'), ('David Fetter'), ('Benjamin Reed');
```

Posts

```
CREATE TABLE message (  
  message_id INTEGER PRIMARY KEY,  
  parent_id INTEGER  
    REFERENCES message(message_id),  
  message_text TEXT NOT NULL,  
  forum_user_id INTEGER  
    NOT NULL REFERENCES forum_users(user_id)  
);
```

Add some posts

```
INSERT INTO message
WITH RECURSIVE m(
  message_id,
  parent_id,
  message_text,
  forum_user_id)
AS (
  VALUES(1, NULL::integer, md5(random()::text), 1)
```

Add some posts

UNION ALL

Add some posts

```
SELECT
  message_id+1,
  CASE
    WHEN random() >= .5 THEN NULL
    ELSE FLOOR(random()*message_id)+1
  END::integer,
  md5(random()::text),
  floor(random() * 10)::integer +1
FROM m
WHERE message_id < 1001
)
SELECT * FROM m;
```

WELL?!?

Patience :)

Find the first post

```
WITH RECURSIVE t1 AS (  
  SELECT  
    /* First message in the thread is the thread ID */  
    message_id AS thread_id,  
    message_id,  
    parent_id,  
    forum_user_id,  
    ARRAY[message_id] AS path  
  FROM message  
  WHERE parent_id IS NULL
```


Find the Next Ones

U N I O N A I L L

Find the Next Ones

```
SELECT
    t1.thread_id,
    m.message_id,
    m.parent_id,
    m.forum_user_id,
    t1.path || m.message_id
FROM message m
JOIN t1 ON
    (t1.message_id = m.parent_id)
),
```

Count Posters in Each Thread

```
t2 AS (  
  SELECT  
    thread_id,  
    forum_user_id,  
    count(*) AS reply_count  
  FROM t1  
  GROUP BY thread_id, forum_user_id  
  ORDER BY thread_id, count(*)  
),
```

Find the Top Posters

```
t3 AS (  
  SELECT thread_id,  
         max(reply_count) AS reply_count  
  FROM t2  
  GROUP BY thread_id  
)
```

Show Them :)

```
SELECT t2.thread_id, f.user_name, t3.reply_count
FROM t2
JOIN t3 USING (thread_id, reply_count)
JOIN forum_users f ON (f.user_id = t2.forum_user_id)
WHERE reply_count > 3
ORDER BY reply_count DESC;
```

Top Posters :)

thread_id	user_name	reply_count
1	Tom Lane	9
1	Gregory Stark	9
82	Magnus Hagander	5
108	Dave Page	4
9	Josh Berkus	4

(5 rows)

OBTW

With CTE and Windowing, SQL is Turing Complete.

Cyclic Tag System

The productions are encoded in the table "p" as follows:

"iter" is the production number;

"rnum" is the index of the bit;

"tag" is the bit value.

This example uses the productions:

110 01 0000

The initial state is encoded in the non-recursive union arm,
in this case just '1'

The $(r.iter \% n)$ subexpression encodes the number of
productions, which can be greater than the size of table "p",
because empty productions are not included in the table.

Cyclic Tag System

Parameters:

the content of "p"

the content of the non-recursive branch

the 3 in $(r.iter \% 3)$

"p" encodes the production rules; the non-recursive branch is the initial state, and the 3 is the number of rules

The result at each level is a bitstring encoded as 1 bit per row, with rnum as the index of the bit number.

At each iteration, bit 0 is removed, the remaining bits shifted up one, and if and only if bit 0 was a 1, the content of the current production rule is appended at the end of the string.

Proof:

Construct a Cyclic Tag System with
CTEs and Windowing.

Proof:

WITH RECURSIVE

p(iter,rnum,tag) AS (

VALUES (0,0,1), (0,1,1), (0,2,0),

(1,0,0), (1,1,1),

(2,0,0), (2,1,0), (2,2,0), (2,3,0)

),

Proof:

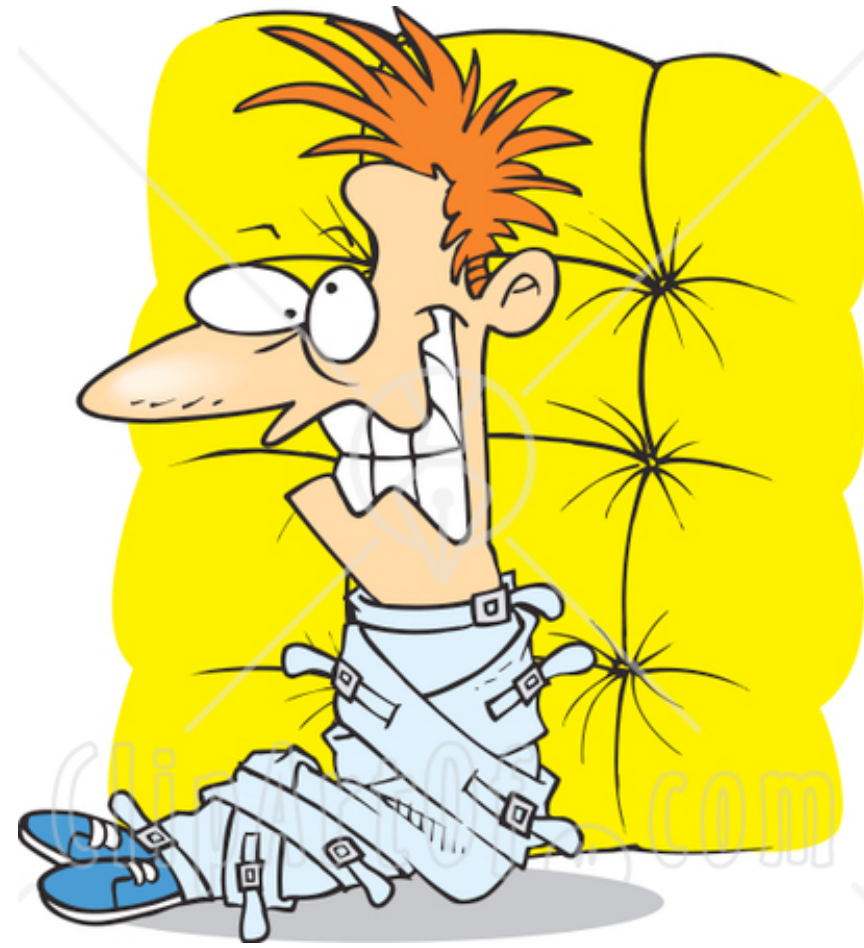
```
r(iter,rnum,tag) AS (  
    VALUES (0,0,1)  
UNION ALL  
    SELECT r.iter+1,  
           CASE  
               WHEN r.rnum=0 THEN p.rnum + max(r.rnum) OVER (  
               ELSE r.rnum-1  
           END,  
           CASE  
               WHEN r.rnum=0 THEN p.tag  
               ELSE r.tag  
           END  
FROM  
    r  
LEFT JOIN p  
    ON (r.rnum=0 and r.tag=1 and p.iter=(r.iter % 3))  
WHERE  
    r.rnum>0  
OR p.iter IS NOT NULL
```

Proof:

```
SELECT iter, rnum, tag  
FROM r  
ORDER BY iter, rnum;
```

Thanks

Andrew (RhodiumToad) Gierth



Questions?

Comments?

Straitjackets?

Thank You!

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